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selection_sort(A[0..n-1])
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k = 0  
while k < n {          (*)  
    i = index j in {k..n-1} for which A[j] is minimal  
    swap A[k] and A[i]  
    k = k + 1  
}
```

Let A' be the original array passed on to the function

Loop invariant: $A[0..k-1]$ contains the smallest k elements of A' in sorted order,
and A is a permutation of A'

Initialization: ($k=0$) $A[0..-1]$ contains the smallest 0 elements of A' - trivially true,
 A is also a permutation of A' (because $A = A'$)

Maintenance: Assume invariant holds at point (*) in the program execution, show that it also holds after the loop body is executed

After the inner loop is executed, i contains an index relating to a smallest element in $A[k..n-1]$. Swapping it with $A[k]$ results in sorted array $A[0..k]$ containing the smallest $k+1$ values of A' , because $A[0..k-1]$ was sorted and $A[k..n-1]$ contained values bigger or equal to $A[k-1]$. Also, A was a permutation of A' . So, when swapping two elements, it still is. k is then incremented and the loop invariant holds once more

Termination: The algorithm stops after exactly n iterations (counting) with $k = n$. Plugging $k=n$ into the loop invariant shows $A[0..n-1]$ contains the smallest elements of A' in sorted order

Q.E.D.